Chapter 5

Least Squares Inverse Eigenvalue Problems

- Overview
- Least Squares IEP
- Theorem
- Geometric Sketch
- Algorithm

Overview

- The information of spectral constraint and the structrual constraint may not be obtained precisely in practice.
- It may occur in practice that one of the two constraints should be enforced more critically than the other.
 - ♦ Without the physical realizability, the physical system simply cannot be built. So structrual constraint is emphasized.
- It may occur that one constraint could be more relaxed than the other.
 - Solution For complex system there is no accurate way to measure the spectrum or to obtain the entire information. Physical uncertainty can be tolerated up to a certain degree.
- When the two constraints cannot be settled simultaneously, the IEP could be formulated in a least squares setting.

- All IEP's discussed so far have a natural generalization to least squares formulation.
- Depends upon which constraint is to be enforced, there are two ways to formulate the least squares IEP.

Least Squares Approximating the Spectrum

• Given:

- \diamond Symmetric matrices $A(d) \in \mathbb{R}^{n \times n}$,
- \diamond Real numbers $\lambda_1^*, \cdots, \lambda_m^*$,
- $\diamond m \le n,$

• Solve:

$$\min_{d \in \mathbb{R}^{l}, \sigma} F(d, \sigma) = \sum_{i=1}^{m} (\lambda_{\sigma_{i}}(d) - \lambda_{i}^{*})^{2},$$
$$= \|\Lambda_{\sigma}(d) - \Lambda_{m}^{*}\|_{F}^{2}.$$

 $\diamond \sigma := \text{a permutation};$ $\diamond \Lambda_m^* := \text{diag}(\lambda_1^*, \dots, \lambda_m^*).$

Least Squares Approximating the Structure

• Let

• Define:

$$\Gamma := \left\{ Q \begin{bmatrix} \Lambda_m^* & 0 \\ 0 & \Lambda_{n-m} \end{bmatrix} Q^T \mid Q \in \mathcal{O}(n), \Lambda_{n-m} \in D_{n-m} \right\}.$$

• Given:

- \diamond Symmetric matrices $A(d) \in \mathbb{R}^{n \times n}$,
- \diamond Real numbers $\lambda_1^*, \cdots, \lambda_m^*$,
- $\diamond m \le n,$
- Solve:

$$\operatorname{dist}(\mathcal{A}(d), \Gamma) = \min_{d \in \mathbb{R}^l, B \in \Gamma} \|A(d) - B\|_F^2.$$

Main Theorem

- At a global minimizer, LSIEP1 and LSIEP2 are equivalent.
 - \diamond Both problems have the same objective values.

$$\min_{d,\sigma} \|\Lambda_{\sigma}(d) - \Lambda_{m}^{*}\|_{F}^{2} = \min_{d,B} \|A(d) - B\|_{F}^{2}.$$

 \diamond If d_1 , σ solve **LSIEP1** and d_2 , B solve **LSIEP2**, then

$$\triangleright d_{1} = d_{2},$$

$$\triangleright B = Q(d_{1}) \begin{bmatrix} \Lambda_{m}^{*} & 0 \\ 0 & \Lambda_{\overline{\sigma}}(d_{1}) \end{bmatrix} Q(d_{1})^{T}.$$

$$\cdot Q(d_{1})^{T} A(d_{1}) Q(d_{1}) = \begin{bmatrix} \Lambda_{\sigma}(d_{1}) & 0 \\ 0 & \Lambda_{\overline{\sigma}}(d_{1}) \end{bmatrix}.$$

$$\cdot \operatorname{diag}[\Lambda_{\sigma}(d_{1}), \Lambda_{\overline{\sigma}}(d_{1})] \text{ has same ordering as } \operatorname{diag}[\Lambda_{m}^{*}, \Lambda_{\overline{\sigma}}(d_{1})]$$

$$\cdot \overline{\sigma} := \{1, ..., n\} - \sigma.$$

 \diamond If so

$$\min_{d,B} \|A(d) - B\|_F^2 = \sum_{i=0}^m |\lambda_{\sigma_i} - \lambda_i^*|^2.$$

One Particular Case

• Consider the special case where

$$A(d) := A_0 + \sum_{i=1}^{l} d_i A_i \in \mathcal{A}(d) \text{ (affine subspace)}$$

- Lift & Project:
 - 1. Find

$$Z^{(k)} := \arg\min_{B\in\Gamma} \|A(d^{(k)}) - B\|_F \quad \text{(Lift)}$$

by the Wielandt-Hoffman Theorem!

2. Find

$$d^{(k+1)} := \arg\min_{d \in \mathbb{R}^l} \|Z^{(k)} - \mathcal{A}(d)\|_F \quad \text{(Project)}$$

by solving a linear system.

• The topology of Γ is not clear. But *Step 1* can be accomplished by considering a "substructure".

$$\mathcal{M}_k := \left\{ Q \begin{bmatrix} \Lambda_m^* & 0\\ 0 & \Lambda_{\overline{\sigma}}(d^{(k)}) \end{bmatrix} Q^T \mid Q \in \mathcal{O}(n) \right\}$$

which is known.

Geometric Sketch



Figure 1: Geometry of life and projection.

• This a descent method!

$$||Z^{(k+1)} - A(d^{(k+1)})||_F \le ||Z^{(k)} - A(d^{(k+1)})||_F$$
$$\le ||Z^{(k)} - A(d^{(k)})||_F$$

Algorithm

- Given $d^{(0)}$.
- For k = 0, 1, 2, ...
 - 1. Compute $Q(d^{(k)})$ and $\Lambda(d^{(k)})$.
 - 2. Find $\sigma = \{\sigma_1, ..., \sigma_m\}$. Form $\Lambda_{\overline{\sigma}}$
 - 3. (lift) Form

$$Z^{(k)} = Q(d^{(k)}) diag(\Lambda^*, \Lambda_{\overline{\sigma}}(d^{(k)})) Q(d^{(k)})^T.$$

4. (**Project**) Compute $d^{(k+1)}$ from

$$\sum_{i=1}^{l} \langle A_i, A_j \rangle d_i^{(k+1)} = \langle Z^{(k)} - A_0, A_j \rangle.$$

5. stop if $||d^{(k+1)} - d^{(k)}|| < \varepsilon$