

Structured Low Rank Approximation

Lecture III: Distance Geometry and Protein Structure

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Syllabus

- Objectives:
 - ◇ To provide some preliminaries.
 - ◇ To treat some mathematics.
 - ◇ To point out some applications.
 - ◇ To describe some algorithms.
- Topics:
 - ◇ Lecture I: Introduction
 - ◇ Lecture II: General Approach
 - ◆ Lecture III: Distance Geometry and Protein Structure
 - ◇ Lecture IV': Singular Value Assignment with Low Rank Matrices
 - ◇ Lecture V: Nonnegative Matrix Factorization
- Assignments:
 - ◇ Application to some real protein data.
 - ◇ Streamline the code to make it portable.

Lecture III

Distance Geometry and Protein Structure

Joint Work with David Chu and Hunter Brown

Outline

- Introduction
 - ◇ Protein Folding Problem
 - ◇ Euclidean Distance Matrix
- Least Squares Formulation
 - ◇ Approximation Problem
 - ◇ Completion Problem
- Analytic Gradient
- Analytic Hessian
- Numerical Examples

Introduction

- Fundamental Problem in Distance Geometry:
 - ◇ Given the distance and chirality constraints which define (our state of knowledge of) a molecule,
 - ◇ Find one or more conformations which satisfy them, or else prove that no such conformation exists.
- Applications:
 - ◇ Molecular conformation problems in chemistry.
 - ◇ Multidimensional scaling in behavioral sciences.
 - ◇ Multivariate analysis in statistics.
 - ◇ Remote exploration and sensing, and antenna array processing.
 - ◇ ...

Protein Folding Problem

- What we know about proteins:
 - ◇ A protein molecule is a connected sequence of amino acid molecules.
 - ◇ There are only twenty amino acids in nature.
 - ▷ Representing each amino acid by a letter from a twenty-letter alphabet, a protein is a string of amino acids linked together like a word.
 - ◇ Most laboratories have the technology to find the ordered sequence of amino acids in a protein.

Real Issue in Protein Folding

- Merely knowing the long linear chains is not enough.
 - ◇ To function properly, the one-dimensional primary amino acid sequence must fold into a particular three-dimensional conformation called its *tertiary configuration*.
 - ◇ This protein then interacts three-dimensionally with other proteins or other groups of molecules called *substrates* in a manner much like a lock and key arrangement.
- The tertiary structure that mediates how a protein functions.
 - ◇ The final folding configuration determines whether a protein acts as an enzyme in chemical reactions or behaves as part of an antibody in the immune system.
 - ◇ How a protein folds determines how the protein works, which ultimately determines how our bodies work.
 - ◇ Wrongly folded proteins do not behave normally, and their abnormal function can lead to disease.
- Understanding how proteins fold into three-dimensional structures given their primary linear sequence thus becomes an incredibly important task.

Then What?

- Literature on the various aspects of protein folding is enormous.
 - ◊ Lots of models using a synthesis of knowledge from biology, physics, chemistry, and mathematics.
 - ◊ No clear models.
- Try to see the three dimensional structure.
 - ◊ Biologists use techniques such as x-ray crystallography or molecular dynamics models to peer through the nano-structure.
 - ◊ Noises in the sensing devices or the imperfection of the models often results in indefinite or incomplete images.
- An imperative task:
 - ◊ Retrieve a possible folding structure that is nearest to the observed but possibly inconsistent or imperfect configuration.

Euclidean Distance Matrix

- Given n particles at locations $\mathbf{p}_1, \dots, \mathbf{p}_n \in \mathbb{R}^m$,
 - ◇ A complete record of relative spacing between any two of the n particles is called a distance matrix.
 - ◇ The matrix $Q(\mathbf{p}_1, \dots, \mathbf{p}_n) = [q_{ij}]$ where

$$q_{ij} = \|\mathbf{p}_i - \mathbf{p}_j\|_2^2, \quad i, j = 1, \dots, n,$$

is called a *Euclidean distance matrix*.

Structures of Euclidean Distance Matrices

- (Schoenberg'35, Young and Householder'38) $Q \in \mathbb{R}^{(n+1) \times (n+1)}$ is a Euclidean distance matrix if and only if $A \in \mathbb{R}^{n \times n}$ defined by

$$a_{ij} := \frac{1}{2}[d_{0i} + d_{0j} - d_{ij}], \quad 1 \leq i, j \leq n$$

is positive semidefinite.

- ◊ $\text{rank}(A)$ is the minimum imbedding dimension, i.e., the lowest dimension of the Euclidean space in which the points reside.
- ◊ Let A be spectrally decomposed as

$$A = U\Lambda U^\top, \tag{1}$$

then columns of $C := \Lambda^{1/2}U^\top$ are coordinates for P_1, \dots, P_n .

- (Gower'82) Q is a distance matrix if and only if $\mathbf{x}^\top Q \mathbf{x} \leq 0$ for all \mathbf{x} such that $\mathbf{x}^\top \mathbf{1} = 0$.
 - ◊ Q is a distance matrix if and only if $(I - \mathbf{1}\mathbf{s}^\top)Q(I - \mathbf{s}\mathbf{1}^\top)$ is negative semidefinite whenever $\mathbf{s}^\top \mathbf{1} = 1$.

Rank Structure

- For any $n \geq m + 2$, the rank of Q is no greater than $m + 2$ and is generically $m + 2$.

◇ Write

$$Q = [\|\mathbf{p}_i - \mathbf{p}_j\|^2] = \sum_{\ell=1}^m [(p_{\ell i} - p_{\ell j})^2].$$

◇ For each ℓ ,

$$Q_\ell = [(p_{\ell i} - p_{\ell j})^2]$$

is a distance matrix for n points $p_{\ell 1}, \dots, p_{\ell n}$ on the real line.

◇ Such a matrix can be denoted in a canonical form,

$$Q_\ell = \begin{bmatrix} 0 & d_1^2 & (d_1 + d_2)^2 & (d_1 + d_2 + d_3)^2 & \dots & & \\ d_1^2 & 0 & d_2^2 & (d_2 + d_3)^2 & & & \\ (d_1 + d_2)^2 & d_2^2 & 0 & d_3^2 & & & \\ \vdots & & & & & & \\ \vdots & & & & & \ddots & \\ & & & & & & 0 \end{bmatrix},$$

where d_1, \dots, d_n denote the 1-dimensional coordinates of these points.

◇ Q_ℓ can further be reduced by eliminations and scalings to

$$\begin{bmatrix} -2 & d_1 & 0 & 0 & 0 & 0 & 0 \\ d_1 & 0 & d_1 + d_2 & d_2 + d_3 & d_3 + d_4 & d_4 + d_5 & d_5 + d_6 \\ 0 & d_1 + d_2 & 0 & 0 & 0 & 0 & 0 \\ 0 & d_2 + d_3 & 0 & 0 & 0 & 0 & 0 \\ 0 & d_3 + d_4 & 0 & 0 & 0 & 0 & 0 \\ 0 & d_4 + d_5 & 0 & 0 & 0 & 0 & 0 \\ 0 & d_5 + d_6 & 0 & 0 & 0 & 0 & 0 \end{bmatrix}.$$

An Inverse Problem

- Once all the inter-particle distances are in hand, it is a simple task to construct the (3-dimensional) geometric structure.
- Not all the inter-particle distances are needed for the construction.
- Many entries in the matrix provide redundant information.
- Rigid body motions, i.e., rotations, translations, and reflections, of particles in space will not alter the distance matrix.
- **Open Question:** Characterize when and what partial information of inter-particle distances will be sufficient for the determination of the entire geometric structure.

Returning to the Scenario

- Let $F \in \mathbb{R}^{n \times n}$ denote an observation of (the squares of) the relative spacing among n particles in \mathbb{R}^3 .
 - ◇ Some of the spacing has not been measured accurately.
 - ▷ $\text{rank}(F) \geq 5$.
 - ▷ Want to retrieve whatever feasible information out of the matrix F so as to infer a realistic conformation.
 - ◇ Some of the spacing is unobservable.
 - ▷ The matrix F is incomplete.
 - ▷ Want to complete the matrix F so that it becomes a distance matrix.
- Find an Euclidean distance matrix $Q \in \mathbb{R}^{n \times n}$ representing the minimal change of F .

Criteria in Approximation

- Criteria used in characterizing the changes affect the way in which the approximation is formulated.
 - ◇ Low rank approximation techniques (Chu et al'02).
 - ▷ A nonnegative and symmetric matrices of rank 5 is not necessarily a distance matrix.
 - ▷ The distance matrix has more structure that is yet to be exploited.
 - ◇ Alternating projection between two geometric entities (Glunt et al'90).
 - ▷ Slow convergence and possible stagnation.
 - ◇ Semi-definite programming (Alafkin'99).
 - ◇ Limited-memory variable-metric continuation techniques (Moré et al'96).
- Least squares formulation with location vectors $\mathbf{p}_1, \dots, \mathbf{p}_n$ as parameter.
 - ◇ The resulting approximation of the imperfect F is guaranteed to be a distance matrix.
 - ◇ Not unique under rotation, translation, or reflection.

Basic Formulation

- Assume initially that all entries of F except the diagonal are to be approximated.
 - ◊ Two special cases later:
 - ▷ The approximation problem where some location vectors of the n particles are known and fixed.
 - ▷ The completion problem where a partial information of spacing among the n particles, but not necessarily from known location vectors, is known and fixed.
- Given $F \in \mathbb{R}^{n \times n}$, minimize the objective function

$$f(\mathbf{p}_1, \dots, \mathbf{p}_n) = \frac{1}{2} \|F - [\langle \mathbf{p}_i - \mathbf{p}_j, \mathbf{p}_i - \mathbf{p}_j \rangle]\|_F^2. \quad (2)$$

- ◊ $\langle \cdot, \cdot \rangle$ = the usual Euclidean inner product.
- ◊ $\|\cdot\|_F$ = the Frobenius matrix norm.
- Our contributions:
 - ◊ The general theory is developed in \mathbb{R}^m .
 - ◊ Offer a highly organized and effective way of computing the derivative information.

Analytic Gradient

- Denote discrepancies

$$d_{ij} = f_{ij} - \langle \mathbf{p}_i - \mathbf{p}_j, \mathbf{p}_i - \mathbf{p}_j \rangle, \quad i, j = 1, \dots, n. \quad (3)$$

- Consider the space $\mathbb{R}^m \times \dots \times \mathbb{R}^m$ equipped with the product topology.

◇ The Fréchet derivative of f at the point $(\mathbf{p}_1, \dots, \mathbf{p}_n)$ acting on an arbitrary n -fold vector $(\mathbf{z}_1, \dots, \mathbf{z}_n) \in \mathbb{R}^m \times \dots \times \mathbb{R}^m$ can be represented as the multi-linear functional

$$f'(\mathbf{p}_1, \dots, \mathbf{p}_n) \cdot (\mathbf{z}_1, \dots, \mathbf{z}_n) = \sum_{k=1}^n \frac{\partial f}{\partial \mathbf{p}_k}(\mathbf{p}_1, \dots, \mathbf{p}_n) \cdot \mathbf{z}_k \quad (4)$$

▷ $A \cdot \mathbf{b}$ denotes the operator action of A on \mathbf{b} .

Partial Gradient

$$\frac{\partial f}{\partial \mathbf{p}_k} = -4 \sum_{\substack{j=1 \\ j \neq k}}^n (\mathbf{p}_k - \mathbf{p}_j) d_{kj}. \quad (5)$$

- Riesz representation:

$$\frac{\partial f}{\partial \mathbf{p}_k}(\mathbf{p}_1, \dots, \mathbf{p}_n) \cdot \mathbf{z}_k = \left\langle \frac{\partial f}{\partial \mathbf{p}_k}(\mathbf{p}_1, \dots, \mathbf{p}_n), \mathbf{z}_k \right\rangle.$$

- Carry on:

$$\left\langle \frac{\partial f}{\partial \mathbf{p}_k}, \mathbf{z}_k \right\rangle = \left\langle \frac{\partial}{\partial \mathbf{p}_k} (F - [\langle \mathbf{p}_i - \mathbf{p}_j, \mathbf{p}_i - \mathbf{p}_j \rangle]) \cdot \mathbf{z}_k, F - [\langle \mathbf{p}_i - \mathbf{p}_j, \mathbf{p}_i - \mathbf{p}_j \rangle] \right\rangle,$$

- Keep going:

$$\begin{aligned} & \frac{\partial}{\partial \mathbf{p}_k} (F - [\langle \mathbf{p}_i - \mathbf{p}_j, \mathbf{p}_i - \mathbf{p}_j \rangle]) \cdot \mathbf{z}_k = \\ & -2 \begin{pmatrix} 0 & \dots & \langle \mathbf{z}_k, \mathbf{p}_k - \mathbf{p}_1 \rangle & & 0 \\ \vdots & & \vdots & & \vdots \\ \langle \mathbf{z}_k, \mathbf{p}_k - \mathbf{p}_1 \rangle & \dots & 0 & \dots & \langle \mathbf{z}_k, \mathbf{p}_k - \mathbf{p}_n \rangle \\ & & \vdots & & \\ 0 & & \langle \mathbf{z}_k, \mathbf{p}_k - \mathbf{p}_n \rangle & & 0 \end{pmatrix}. \end{aligned}$$

- Finally,

$$\frac{\partial f}{\partial \mathbf{p}_k}(\mathbf{p}_1, \dots, \mathbf{p}_n) \cdot \mathbf{z}_k = \left\langle \mathbf{z}_k, -4 \sum_{\substack{j=1 \\ j \neq k}}^n (\mathbf{p}_k - \mathbf{p}_j) d_{kj} \right\rangle$$

Gradient of f

- Using the product topology again,

$$\nabla f(\mathbf{p}_1, \dots, \mathbf{p}_n) = \left(\frac{\partial f}{\partial \mathbf{p}_1}, \dots, \frac{\partial f}{\partial \mathbf{p}_n} \right) \in \mathbb{R}^m \times \dots \times \mathbb{R}^m. \quad (6)$$

- Any numerical method utilizing the gradient information can now be employed to solve the approximation problem.
- Gradient flow:

$$\frac{d\mathbf{p}_k}{dt} = 4 \sum_{\substack{j=1 \\ j \neq k}}^n (\mathbf{p}_k - \mathbf{p}_j) (f_{kj} - \langle \mathbf{p}_k - \mathbf{p}_j, \mathbf{p}_k - \mathbf{p}_j \rangle), \quad k = 1, \dots, n$$

moves in the steepest descent direction to reduce the values for the objective function f .

- ◇ A descent flow bounded in a neighborhood of F .
- ◇ Limit point must exist.
- ◇ A (local) least squares approximation of the given F .

Analytic Hessian

- Will be more effective for finding critical points if the second-derivative information of f is available.

- Consider

$$g : \mathbb{R}^m \times \dots \times \mathbb{R}^m \rightarrow \mathbb{R}^m \times \dots \times \mathbb{R}^m.$$

$$\diamond g = \nabla f = (g_1, \dots, g_n).$$

$$g_k(\mathbf{p}_1, \dots, \mathbf{p}_n) = \frac{\partial f}{\partial \mathbf{p}_k}(\mathbf{p}_1, \dots, \mathbf{p}_n). \tag{7}$$

- The Jacobian matrix of g_k constitutes precisely the k -th row block of the Hessian of f .

$$\diamond \text{The Hessian can be calculated block by block.}$$

Partial Jacobian

$$\frac{\partial g_k}{\partial \mathbf{p}_i} = \begin{cases} \sum_{j=1, j \neq k}^n [-4d_{kj}I_m + 8(\mathbf{p}_k - \mathbf{p}_j)(\mathbf{p}_k - \mathbf{p}_j)^T], & \text{if } i = k; \\ 4d_{ki}I_m - 8(\mathbf{p}_k - \mathbf{p}_i)(\mathbf{p}_k - \mathbf{p}_i)^T, & \text{if } i \neq k. \end{cases} \quad (8)$$

- Under the product topology,

$$g'_k(\mathbf{p}_1, \dots, \mathbf{p}_n) \cdot (\mathbf{w}_1, \dots, \mathbf{w}_n) = \sum_{i=1}^n \frac{\partial g_k}{\partial \mathbf{p}_i} \cdot \mathbf{w}_i.$$

- If $i = k$,

$$\begin{aligned} \frac{\partial g_k}{\partial \mathbf{p}_k} \cdot \mathbf{w}_k &= -4 \sum_{\substack{j=1 \\ j \neq k}}^n [\mathbf{w}_k d_{kj} - 2(\mathbf{p}_k - \mathbf{p}_j) \langle \mathbf{w}_k, \mathbf{p}_k - \mathbf{p}_j \rangle] \\ &= \sum_{\substack{j=1 \\ j \neq k}}^n [-4d_{kj}I_m + 8(\mathbf{p}_k - \mathbf{p}_j)(\mathbf{p}_k - \mathbf{p}_j)^T] \mathbf{w}_k. \end{aligned}$$

- If $i \neq k$,

$$\begin{aligned} \frac{\partial g_k}{\partial \mathbf{p}_i} \cdot \mathbf{w}_i &= -4[-\mathbf{w}_i d_{ki} - 2(\mathbf{p}_k - \mathbf{p}_i) \langle -\mathbf{w}_i, \mathbf{p}_k - \mathbf{p}_i \rangle] \\ &= [4d_{ki}I_m - 8(\mathbf{p}_k - \mathbf{p}_i)(\mathbf{p}_k - \mathbf{p}_i)^T] \mathbf{w}_i. \end{aligned}$$

Modification

- Usually some additional information is available in the geometric conformation.
 - ◇ The x-ray crystal structure of each of the twenty amino acids in nature is known.
 - ◇ Most of the amino acid sequences of our proteins are also known.
 - ◇ Once a certain amino acid is known to present in the protein, a certain block of the matrix F is already predetermined and fixed.
- Least squares formulations should be modified accordingly to reflect this fact.
 - ◇ Derivative information is available in block form.
 - ◇ Convenient for the overall process of assembling essential gradient and the Hessian.

Approximation with Partially Fixed Locations

- Any rotations, translations, or reflections of a given conformation will produce the same relative spacing and hence the same distance matrix.
 - ◊ To shun rotation, translation, or reflection, $m + 1$ positions of these n particles in the embedding space \mathbb{R}^m must be specified and bound as reference points.
 - ◊ Some additional location vectors among $\mathbf{p}_1, \dots, \mathbf{p}_n$ might be known and fixed beforehand.
- Let \mathbf{q} denote the indices of known location vectors.
 - ◊ Entries f_{ij} of F where both $i, j \in \mathbf{q}$ correspond to the spacing among these known location vectors.
 - ◊ These entries of F should be exact and kept constant.
 - ◊ Derivatives at these points should be zero.
 - ◊ Keep these known position vectors invariant by simply nullifying any derivative information at the corresponding blocks.

Completion with Partially Fixed F

- The matrix F represents a partially specified distance matrix.
 - ◊ The completion problem is different from the approximation problem.
 - ◊ The specified entries in F do not necessarily correspond to any known location vectors, but is required to remain the same throughout the whole matrix completion process.
 - ◊ The specified entries in F must be consistent by themselves to begin with.
 - ◊ **Open Question:** Determine whether the specified entries in F are consistent so that F indeed can be completed as a distance matrix (Trosett'97).

- Let Δ denote the index set of those specified entries of F .

$$\Delta := \{(i, j) \in \mathbb{Z} \times \mathbb{Z} \mid d_{ij} = 0\}.$$

- Minimize the same objective function $f(\mathbf{p}_1, \dots, \mathbf{p}_n)$ as before, subject to additional equality constraints

$$\langle \mathbf{p}_i - \mathbf{p}_j, \mathbf{p}_i - \mathbf{p}_j \rangle = f_{ij} \quad \text{for all } (i, j) \in \Delta. \tag{9}$$

- ◊ The completion problem is cast as an equality constrained optimization problem.

Example 1

- A knot \mathbb{R}^3 .

$$\begin{cases} x = -10 \cos(t) - 2 \cos(5t) + 15 \sin(2t); \\ y = -15 \cos(2t) + 10 \sin(t) - 2 \sin(5t); \\ z = 10 \cos(3t); \end{cases} \quad 0 \leq t \leq 2\pi$$

- Simulation:

- ◇ Represent the knot by n discrete points.
- ◇ Use the points to define the true $n \times n$ distance matrix $Q = [q_{ij}]$.
- ◇ Let \mathbf{q} denote the indices of known location vectors.
 - ▷ \mathbf{q} must contain at least four points.
- ◇ Perturb Q to simulate F .

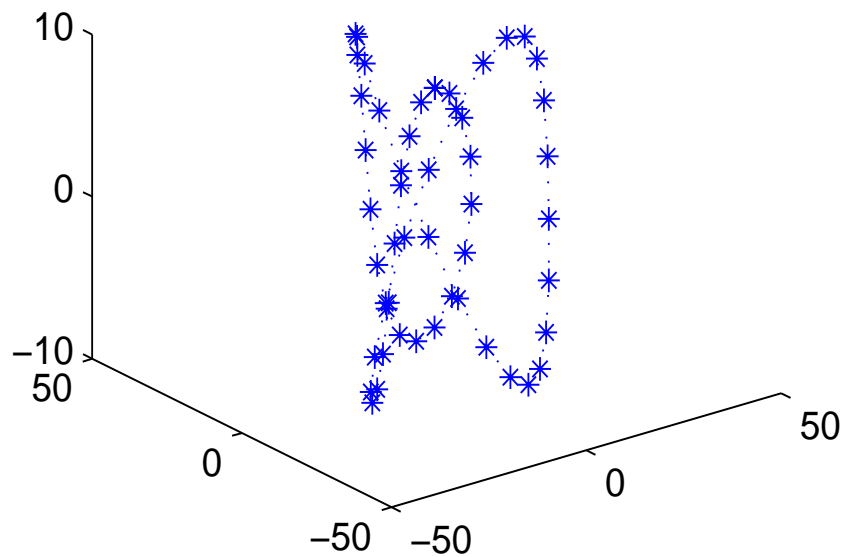
$$f_{ij} = f_{ji} = \begin{cases} 0, & \text{if } i = j, \\ q_{ij}, & \text{if both } i \text{ and } j \text{ are in } \mathbf{q}, \\ (\sqrt{q_{ij}} + \sigma * \text{randn}(1))^2, & \text{if either } i \text{ or } j \text{ is not in } \mathbf{q}. \end{cases}$$

- ▷ σ is an indicator of how far F is away from a distance matrix.
- ▷ $\sigma = 2$ would bury Q in a significant amount of random noise.
- ◇ Use existing routine FMINUNC in the Optimization Toolbox, Version 2.2, of MATLAB.
- ◇ Keep location vectors in by \mathbf{q} fixed throughout the iteration.
- ◇ Perturb each entry of the true location vectors by an additive noise with uniform distribution over $[-20,20]$.

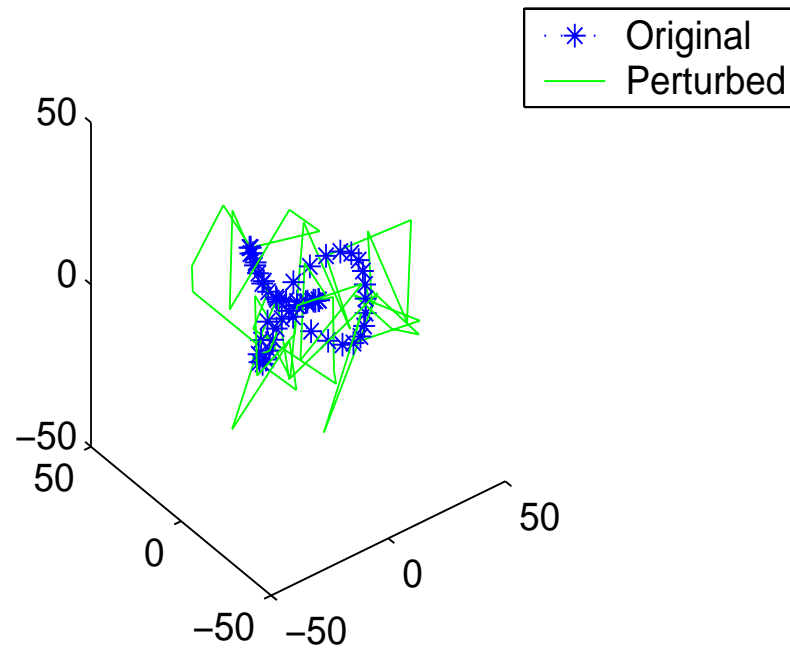
- Numerical results:

- ◇ Objective values have been reduced from an order of 10^9 to 10^6 , indicating that F is far from being a distance matrix.
- ◇ Remarkable likeness between the recovered and the original configurations.

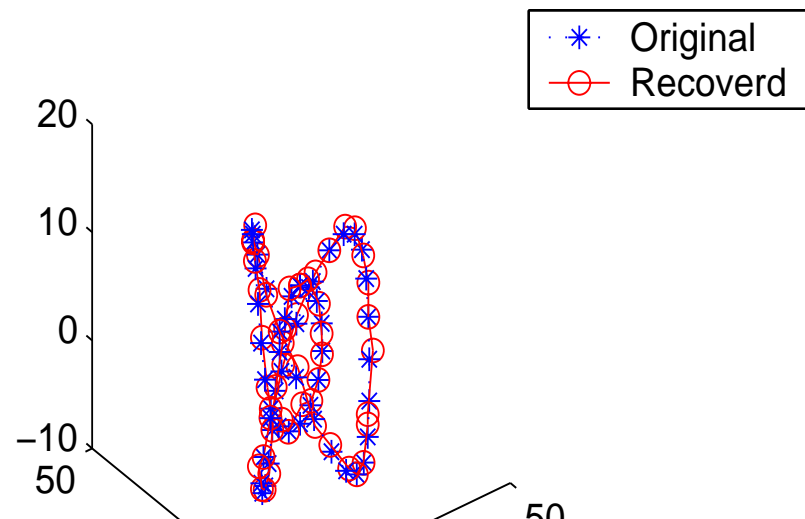
Original Configuration



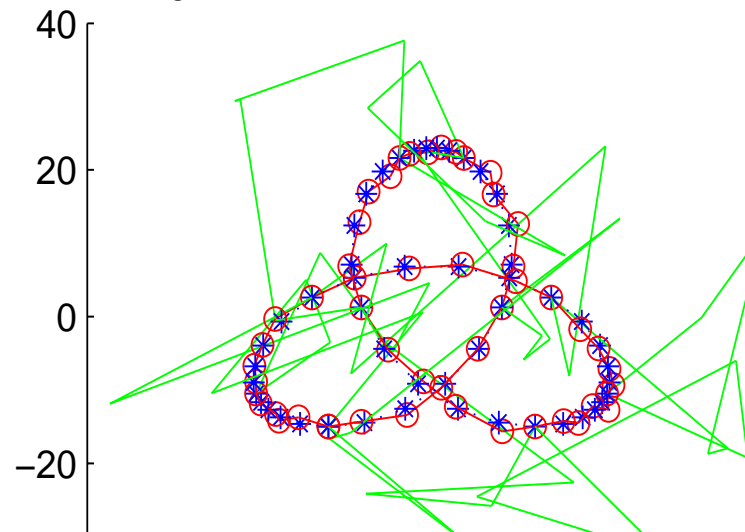
Original and Perturbed Configurations



Original and Recovered Configurations



2D View of Original, Perturbed, and Recovered Configurations



Example 2.

- A helix

$$\begin{cases} x = 4 \cos(3t); \\ y = 4 \sin(3t); & 0 \leq t \leq 2\pi; \\ z = 2t. \end{cases}$$

- Numerical results:

- ◇ $n = 101$ and $\mathbf{q} = 1 : 5 : 101$.

- ▷ Not a smooth solution.

- ▷ Recapture the helix feature from a fairly deviate initial guess and a fraudulent F .

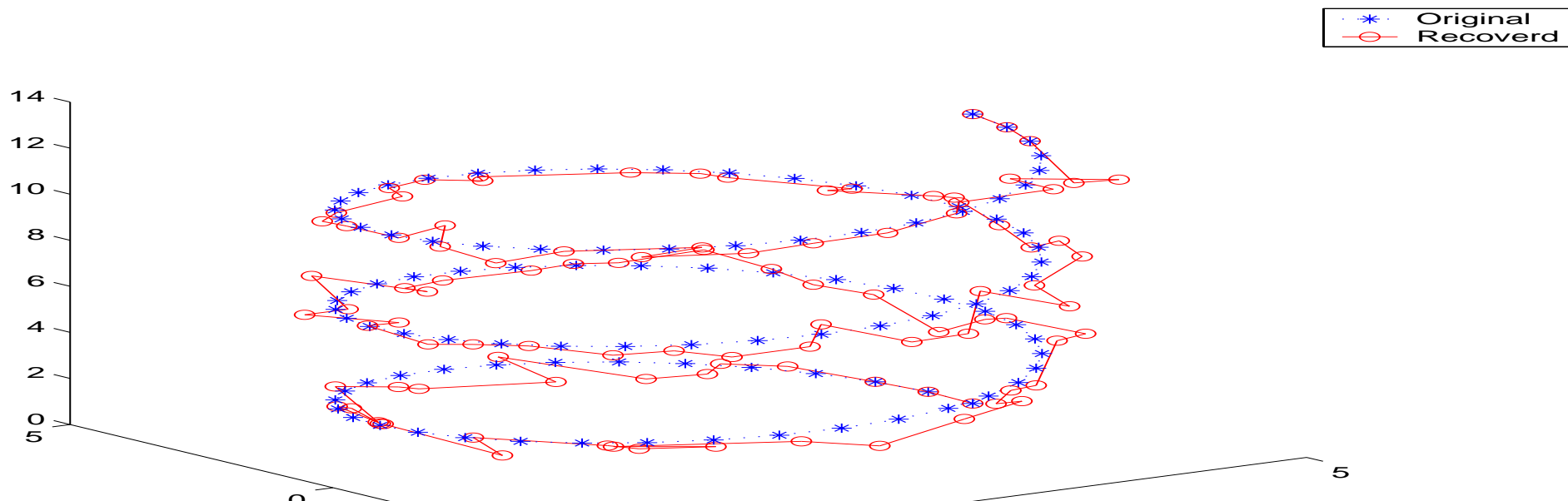
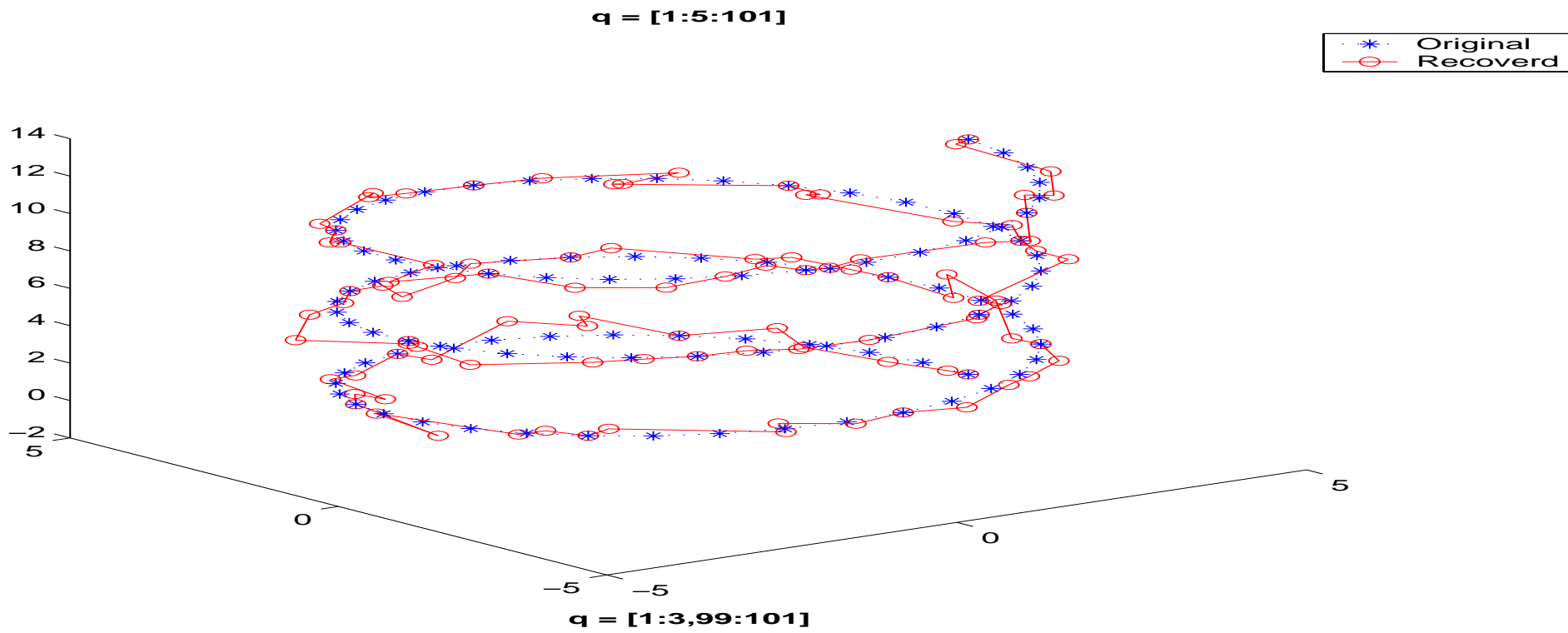
- ▷ Pre-shaped?

- ◇ $\mathbf{q} = [1 : 3, 99 : 101]$.

- ▷ Only the first and the last three location vectors are fixed.

- ▷ Deviate from the helix further.

- ▷ Spirality is evident.



Example 3

- A 2-dimensional pretzel:

$$\begin{cases} x = t + 3 \sin(2t); \\ y = t + 2 \sin(3t); \end{cases} \quad -2\pi \leq t \leq 2\pi.$$

- ◇ Has many critical turns which are close to each other.
- ◇ A more careful observation of the spacing among the location vectors is necessary.
 - ▷ To affect this scenario, we assume a smaller standard deviation $\sigma = 1$ of noise in the simulation.

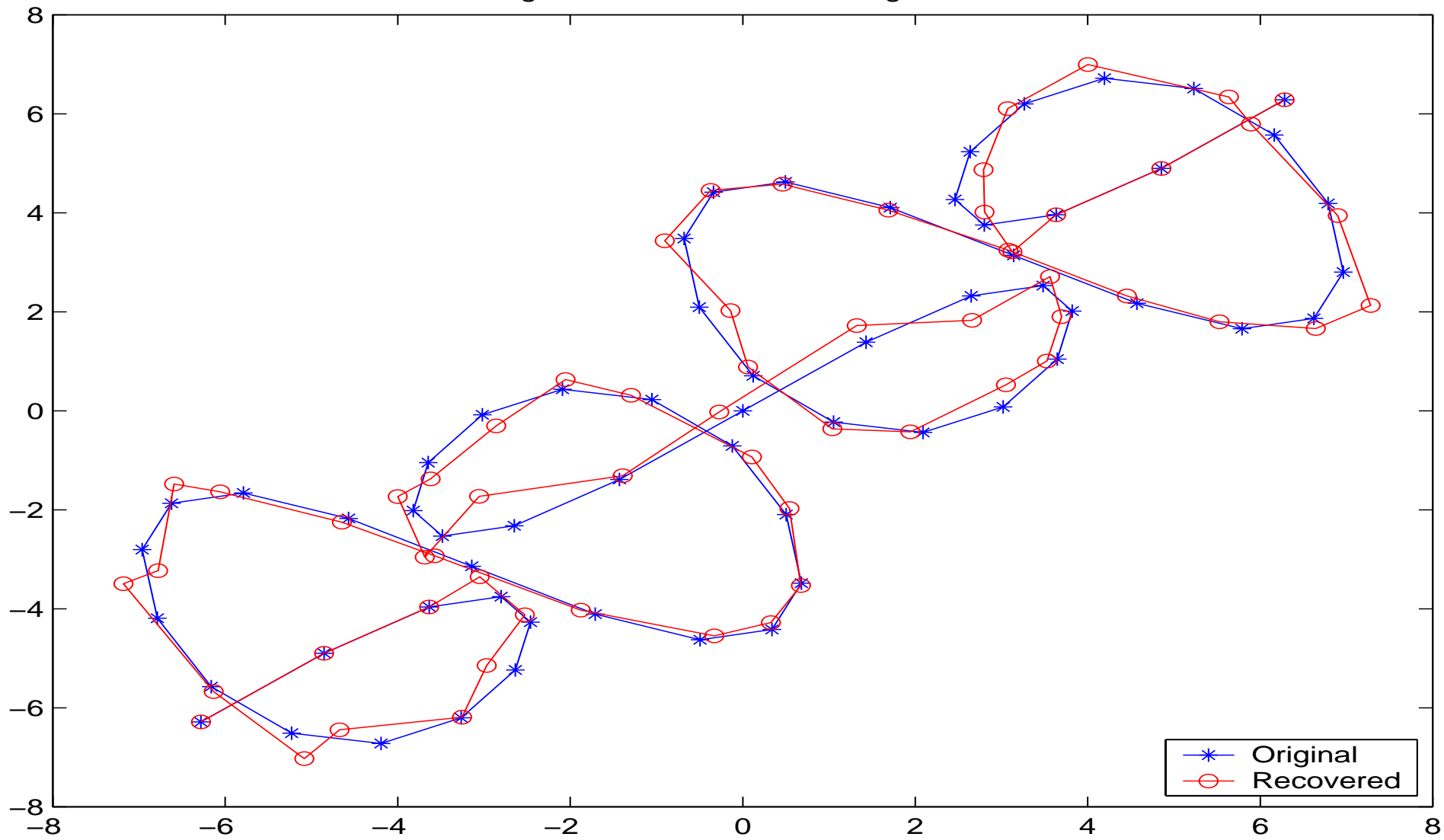
- Simulation:

- ◇ $n = 61$.
- ◇ $\mathbf{q} = [1 : 3, 59 : 61]$.
- ◇ Moderately wild initial guess.

- Numerical results:

- ◇ Able to pick up the folds.
- ◇ Smaller σ or more known position would ease the difficulties.

Original and Recovered Configurations



Example 4

- Complete the 6×6 partially specified matrix

$$F = \begin{bmatrix} 0 & 3 & 4 & 3 & 4 & 3 \\ 3 & 0 & 1 & x_1 & 5 & x_3 \\ 4 & 1 & 0 & 5 & x_2 & 5 \\ 3 & x_1 & 5 & 0 & 1 & x_4 \\ 4 & 5 & x_2 & 1 & 0 & 5 \\ 3 & x_3 & 5 & x_4 & 5 & 0 \end{bmatrix}$$

with values of x_i , $i = 1, \dots, 4$.

- ◇ Look for six location vectors in \mathbb{R}^m .
 - ◇ Not sure about the dimension m of the embedding space.
 - ◇ The first 3×3 principal submatrix is completely known, suggesting that three location vectors could have been self-determined.
- If $m = 2$,
 - ◇ Use $\mathbf{p}_1 = (0, 0)$, $\mathbf{p}_2 = (\sqrt{3}, 0)$, and $\mathbf{p}_3 = (\sqrt{3}, 1)$ as reference points in \mathbb{R}^2 .
 - ◇ Eight equality constraints in the form of (9) for the remaining three location vectors.
 - ◇ Unknown location vectors \mathbf{p}_j , $j = 4, 5, 6$, constitute only six unknowns.
 - ◇ No feasible solution at all in \mathbb{R}^2 .

- If $m = 3$,
 - ◊ Embed \mathbf{p}_i , $i = 1, 2, 3$, in \mathbb{R}^3 by adding zeros to their third components.
 - ◊ Employ an existing routine FMINCON in the Optimization Toolbox of MATLAB by supplying the equality constraints (9).
- Numerical results.

x_1	x_2	x_3	x_4
6.6883	3.9512	2.0187	7.3255
1.7434	9.1772	2.2007	2.2006
2.2800	9.4157	2.3913	4.7487
2.7971	5.7203	7.2315	7.2315
2.2723	9.4208	2.3964	4.7398

Table 1: Examples of entries for completed distance matrix.

\mathbf{p}_1	\mathbf{p}_2	\mathbf{p}_3	\mathbf{p}_4	\mathbf{p}_5	\mathbf{p}_6
0	1.7321	1.7321	0.9014	0.5774	1.1359
0	0	1.0000	-0.5613	-0.4223	-0.9675
0	0	0	1.3335	1.7980	-0.2711
0	0	0	-0.3067	0.5057	0.8367

Table 2: Example of location vectors in \mathbb{R}^4 for completed distance matrix.

Conclusion

- There are several algorithms available for solving the distance matrix approximation problem.
 - ◊ We cast the problem under a least squares approximation in terms of the location vectors directly and propose using conventional large-scale optimization techniques instead.
- We manage the resulting complexity by organizing the gradient and Hessian information in block forms.
 - ◊ The matrix calculus makes it particularly easy to assemble the derivatives for existing software packages when some locations vectors are known and fixed.
- Numerical experiments seem to suggest that the conventional methods are efficient and robust in the reconstruction.