Gaussian Elimination

Consider the Consider the system

$$
\begin{cases}\n a_{11}x_1 + a_{12}x_2 + \ldots + a_{1n}x_n = b_1 \\
 a_{12}x_1 + a_{22}x_2 + \ldots + a_{2n}x_n = b_2 \\
 a_{n1}x_1 + a_{n2}x_2 + \ldots + a_{nn}x_n = b_n,\n\end{cases}
$$

Gauss elimination method consists in

- Using the rst equation assuming a to eliminate all x terms from the second equation on
- $-$ terms from \sim terms from the new second equation to all x-resonance all \sim \sim the second experiments from \sim the third equation on
- , we have a new system and we obtain a new system of the end of the system of the system of the system of the s

the contract of the contract of

$$
\begin{cases} \overline{a}_{11}x_1 + \overline{a}_{12}x_2 + \ldots + \overline{a}_{1n}x_n = \overline{b}_1 \\ \overline{a}_{22}x_2 + \ldots + \overline{a}_{2n}x_n = \overline{b}_2 \\ \vdots \\ \overline{a}_{nn}x_n = \overline{b}_n \end{cases}
$$

• The elimination process can be summarized as follows:

```
if a construction of the construction of t
                        quit
          else
                      wj -
 akj
                      eta -
 aik
akk
                      aik
 -
 eta
                      <u>for a formulation of the set of th</u>
                                aij
 -
 aij
  eta 
 wj
                     end
          end
end
```
- The above algorithm computes the factorization A LU
- The strictly lower triangular portion of A is overwritten by ele ments of L whose diagonal elements are constantly 1.
- The upper triangular portion of A is overwritten by elements of
- T The algorithm will term algorithm will term will term α and α zero α zero α α pivot element.
- In general, after $k-1$ steps, we should have constructed the matrix

$$
A^{(k)} := \left[\begin{array}{cccc} a^{(1)}_{11} & a^{(1)}_{12} & & \ldots & a^{(1)}_{1n} \\ 0 & a^{(2)}_{22} & & a^{2}_{2n} \\ & & \ddots & & \vdots \\ 0 & & 0 & a^{(k)}_{kk} & \ldots & a^{(k)}_{kn} \\ 0 & & 0 & a^{(k)}_{nk} & \ldots & a^{(k)}_{nn} \end{array}\right]
$$

 \Diamond Assuming $a_{kk}^{r'} \neq 0$, we define the multipliers $m_{ik} := a_{ik}^{r'}/a_{kk}^{r'}$ for i k - ---n

- \Diamond Then we generate a_{ij}^{s+1} := $a_{ij}^{s+1} m_{ik} a_{ki}^{s+1}$ for all $i, j = k+1, \ldots, n$.
- In do so the earlier k rows are left unchanged and zeros are introduced into column k below the diagonal element.
- The elimination process can be described in terms of elementary matrix operations
	-

$$
E_{ji}(p) := \begin{bmatrix} 1 & & & & \\ & \ddots & & & \\ j-\text{th row} & \longrightarrow & p & \\ & & p & & \\ 0 & & & 1 \end{bmatrix}
$$

 $i-\text{th column}$

- $\triangleright E_{ji}(p)A$ amounts to multiplying the *i*-th row of A by the scalar p and then adding the result to the j-th row of A .
- \triangleright det $E_{ji}(p) = 1$. φ $(E_{ji}(p))^{-1} = E_{ji}(-p).$
- \diamond The transformation from A^{\succ} to A^{\succ} can be summarized as

$$
E_{n1}(-m_{n1})\dots E_{21}(-m_{21})A^{(1)}=A^{(2)}.
$$

 \diamond -Likewise, A^\vee can be transformed to A^\vee chrough

$$
E_{n2}(-m_{n2})\dots E_{32}(-m_{32})A^{(2)}=A^{(3)}.
$$

- In the end we have the end of the end of the

$$
\underbrace{E_{n,n-1}E_{n-1,n-2}\ldots E_{n2}\ldots E_{32}E_{n1}\ldots E_{21}}_{L^{-1}}A^{(1)}=A^{(n)},
$$

where $A^{\scriptscriptstyle(w)}$ is an upper triangular matrix, and L^{-1} , being the product of a sequence of lower triangular matrices, is a nonsingular lower triangular matrix

 \diamond it remains to check out what L $\,$ $\,$ is.